# Approfondimento: <br> Sincronizzazione 

Architettura dei Calcolatori
Prof. Andrea Marongiu
andrea.marongiu@unimore.it
Anno accademico 2018/19

## Motivation: "Too much milk"

- Example: People need to coordinate:

| Time | Person A | Person B |
| :--- | :--- | :--- |
| $3: 00$ | Look in Fridge. Out of milk |  |
| $3: 05$ | Leave for store |  |
| $3: 10$ | Arrive at store | Look in Fridge. Out of milk |
| $3: 15$ | Buy milk | Leave for store |
| $3: 20$ | Arrive home, put milk away | Arrive at store |
| $3: 25$ |  | Buy milk |
| $3: 30$ |  | Arrive home, put milk away |

## Shared Memory Synchronization

- Program is a collection of processors (or threads of control).
- Each processor/thread has a set of private variables (e.g., local stack variables)
- Also a set of shared variables, (e.g., static variables, or global heap).
- Processors communicate implicitly by writing and reading shared variables.
- Processors coordinate by synchronizing on shared variables



## Shared Memory code for computing a

 sum

- Problem is a race condition on variable $\boldsymbol{s}$ in the program
- A race condition or data race occurs when:
- two processors (or two threads) access the same variable, and at least one does a write.
- The accesses are concurrent (not synchronized) so they could happen simultaneously


## Shared Memory code for computing a

 sum
static int $\mathrm{s}=\mathbf{0}$;

```
Thread 1
    ....
    compute f([A[i]) and put in reg0
    reg1 = s
    reg1 = reg1 + reg0
    s = reg1
```


## Thread 2

...
compute $f([A[i])$ and put in reg0 reg1 $=\mathrm{s}$
reg1 = reg1 + reg0
$\mathrm{s}=\mathrm{reg} 1$

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int s = 0;

```
Thread 1
    *
    compute f([A[i]) and put in reg0
    reg1 = s
    reg1 = reg1 + reg0
    s = reg1
```

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int s = 0;

```
Thread }
    ....
    compute f([A[i]) and put in reg0
    reg1 = s
    reg1 = reg1 + reg0
    s = reg1
```

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int s = 0;

| Thread 1 |  |
| :--- | :--- |
| $\ldots .$. | 9 |
| compute $f([A[i])$ and put in reg0 | 0 |
| reg1 = s |  |
| reg1 = reg1 + reg0 |  |
| $s=$ reg1 |  |
| $\ldots$ |  |

```
Thread 2
    ...
    compute f([A[i]) and put in reg0
    reg1 = s
    reg1 = reg1 + reg0
    s = reg1
```

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int $\mathrm{s}=\mathbf{0}$;

| Thread 1 |  |
| :--- | :--- |
| $\ldots .$. | 9 |
| compute $f([A[i])$ and put in reg0 | 0 |
| reg1 = s |  |
| reg1 = reg1 + reg0 |  |
| $s=$ reg1 |  |
| $\ldots$ |  |

```
Thread 2
    ...
    compute f([A[i]) and put in reg025
    reg1 = s
    reg1 = reg1 + reg0
    s = reg1
```

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int $\mathrm{s}=\mathbf{0}$;

| Thread 1 |  |
| :--- | :--- |
| $\ldots .$. | 9 |
| compute $f([A[i])$ and put in reg0 | 0 |
| reg1 = s |  |
| reg1 = reg1 + reg0 |  |
| $s=$ reg1 |  |
| $\ldots$ |  |

```
Thread 2
    ...
    compute f([A[i]) and put in reg025
    reg1 = s 0
    reg1 = reg1 + reg0 25
    s = reg1
```

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a

 sum
static int s = 0;

| Thread 1 |  |
| :--- | :--- |
| $\ldots .$. | 9 |
| compute $f([A[i])$ and put in reg0 | 0 |
| reg1 $=s$ |  |
| reg1 = reg1 + reg0 |  |
| $s=r e g 1$ |  |
| $\ldots$ |  |

```
Thread 2
    ...
    compute f([A[i]) and put in reg025
    reg1 = s 0
    reg1 = reg1 + reg0 25
    s = reg1
- Assume \(\boldsymbol{A}=[3,5], \boldsymbol{f}\) is the square function, and \(\boldsymbol{s}=0\) initially
- For this program to work, \(\boldsymbol{s}\) should be 34 at the end

\section*{Shared Memory code for computing a sum}

static int s = 0;
\begin{tabular}{|ll|}
\hline Thread 1 & \\
\(\ldots .\). & 9 \\
compute \(f([A[i])\) and put in reg0 & 0 \\
reg1 = s & 9 \\
reg1 = reg1 + reg0 & \\
s = reg1 & \\
\(\ldots\) &
\end{tabular}
```

Thread 2
...
compute f([A[i]) and put in reg025
reg1 = s 0
reg1 = reg1 + reg0 25
s = reg1

- Assume $\boldsymbol{A}=[3,5], \boldsymbol{f}$ is the square function, and $\boldsymbol{s}=0$ initially
- For this program to work, $\boldsymbol{s}$ should be 34 at the end


## Shared Memory code for computing a sum


static int s = 0;

| Thread 1 |  |
| :--- | :--- |
| $\ldots .$. | 9 |
| compute $f([A[i])$ and put in reg0 | 0 |
| reg1 = s | 9 |
| reg1 = reg1 + reg0 | 9 |
| s = reg1 |  |
| $\ldots$ |  |

```
Thread 2
...
compute f([A[i]) and put in reg025
    reg1 = s 0
    reg1 = reg1 + reg0 25
    s = reg1
- Assume \(\boldsymbol{A}=[3,5], \boldsymbol{f}\) is the square function, and \(\boldsymbol{s}=0\) initially
- For this program to work, \(\boldsymbol{s}\) should be 34 at the end

\section*{Shared Memory code for computing a} sum

static int s = 0;
\begin{tabular}{|ll|}
\hline Thread 1 & \\
\(\ldots .\). & 9 \\
compute \(f([A[i])\) and put in reg0 & 0 \\
reg1 = s & 9 \\
reg1 = reg1 + reg0 & 9 \\
s = reg1 & \\
\(\ldots\) &
\end{tabular}
```

Thread 2
...
compute f([A[i]) and put in reg025
reg1 = s 0
reg1 = reg1 + reg0 25
s = reg1 25

```
- Assume \(\boldsymbol{A}=[3,5], \boldsymbol{f}\) is the square function, and \(\boldsymbol{s}=0\) initially
- For this program to work, \(\boldsymbol{s}\) should be 34 at the end
- but it may be 34,9 , or 25

\section*{Shared Memory code for computing a} sum
static int \(s=0\);
```

Thread 1
local_s1= 0
for i=0,n/2-1
local_s1 = local_s1 + sqr(A[i])
s = s + local_s1

```

Thread 2
local_s2 = 0
for \(i=n / 2, n-1\)
local_s2= local_s2 + sqr(A[i])
\(\mathrm{s}=\mathrm{s}+\) local_s2

\section*{Shared Memory code for computing a} sum
static int \(\mathrm{s}=0\);


\section*{Atomic Operations}
- To understand a concurrent program, we need to know what the indivisible operations are!
- Atomic Operation: an operation that always runs to completion or not at all
- It is indivisible: it cannot be stopped in the middle and state cannot be modified by someone else in the middle
- Fundamental building block - if no atomic operations, then have no way for threads to work together
- On most machines, memory references and assignments (i.e. loads and stores) of words are atomic

\section*{Definitions}
- Synchronization: using atomic operations to ensure cooperation between threads
- For now, only loads and stores are atomic
- hard to build anything useful with only reads and writes
- Mutual Exclusion: ensuring that only one thread does a particular thing at a time
- One thread excludes the other while doing its task
- Critical Section: piece of code that only one thread can execute at once
- Critical section and mutual exclusion are two ways of describing the same thing
- Critical section defines sharing granularity

\section*{More Definitions}
- Lock: prevents someone from doing something
- Lock before entering critical section and before accessing shared data
- Unlock when leaving, after accessing shared data
- Wait if locked
- Important idea: all synchronization involves waiting
- Example: fix the milk problem by putting a lock on refrigerator
- Lock it and take key if you are going to go buy milk


\section*{Shared Memory code for computing a} sum
static int \(s=0\);
```

Thread 1
local_s1= 0
for i=0,n/2-1
local_s1 = local_s1 + sqr(A[i])
s = s + local_s1

```

Thread 2
local_s2 = 0
for \(i=n / 2, n-1\)
local_s2= local_s2 + sqr(A[i])
\(\mathrm{s}=\mathrm{s}+\) local_s2

\section*{Shared Memory code for computing a} sum
static int \(\mathrm{s}=0\);


\section*{Shared Memory code for computing a} sum
```

static int s=0;
static lock Ik;

```
```

Thread 1
local_s1= 0
for i=0,n/2-1
local_s1 = local_s1 + sqr(A[i])
lock(lk);
s = s + local_s1
unlock(lk);

```
```

Thread 2
local_s2 = 0
for i = n/2, n-1
local_s2= local_s2 + sqr(A[i])
lock(lk);
s = s +local_s2
unlock(lk);

```

\section*{How to implement locks?}
- Need HW support for atomic instructions
- RISCV uses two HW primitives
- Load reserved
- Store conditional
- (see discussion in Chapter 2: Instructions, language of the computer, slides 61-63)```

