#### Approfondimento: Sincronizzazione

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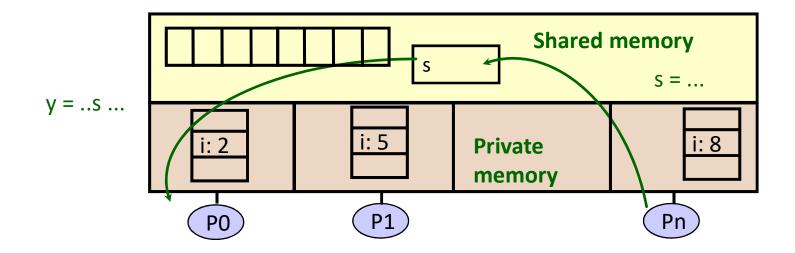


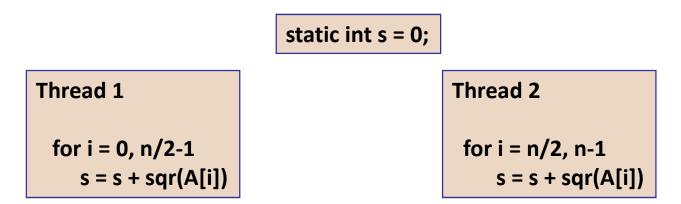
• Example: People need to coordinate:

Time	Person A	Person B
3:00	Look in Fridge. Out of milk	
3:05	Leave for store	
3:10	Arrive at store	Look in Fridge. Out of milk
3:15	Buy milk	Leave for store
3:20	Arrive home, put milk away	Arrive at store
3:25		Buy milk
3:30		Arrive home, put milk away

#### Shared Memory Synchronization

- Program is a collection of **processors** (or *threads* of control).
- Each processor/thread has a set of **private variables** (e.g., local stack variables)
- Also a set of **shared variables**, (e.g., static variables, or global heap).
  - Processors communicate **implicitly** by writing and reading shared variables.
  - Processors coordinate by **synchronizing** on shared variables





- Problem is a race condition on variable s in the program
- A race condition or data race occurs when:
  - two processors (or two threads) access the same variable, and at least one does a write.
  - The accesses are concurrent (not synchronized) so they could happen simultaneously

A 3 5 f = square static int s =	= 0;
Thread 1	Thread 2
compute f([A[i]) and put in reg0	compute f([A[i]) and put in reg0
reg1 = s	reg1 = s
reg1 = reg1 + reg0	reg1 = reg1 + reg0
s = reg1	s = reg1

- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, s should be 34 at the end

f = square Α 3 5 static int s = 0; Thread 1 Thread 2 ... 9 compute f([A[i]) and put in reg0 compute f([A[i]) and put in reg0 reg1 = sreg1 = sreg1 = reg1 + reg0reg1 = reg1 + reg0s = reg1s = reg1... ...

- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, **s** should be 34 at the end

A 3 5 <b>f = square</b> static int s =	= 0;
Thread 1	Thread 2
compute f([A[i]) and put in reg0 9	compute f([A[i]) and put in reg0 25
reg1 = s	reg1 = s
reg1 = reg1 + reg0	reg1 = reg1 + reg0
s = reg1	s = reg1

- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, s should be 34 at the end

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- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, s should be 34 at the end

A 3 5 f = square static int s	= 0;
Thread 1	Thread 2
 compute f([A[i]) and put in reg0 9	 compute f([A[i]) and put in reg0 25
reg1 = s 0	$reg1 = s \qquad 0$
reg1 = reg1 + reg0 s = reg1	reg1 = reg1 + reg0       25         s = reg1       25

- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, s should be 34 at the end

A 3 5 f = square sta	<mark>atic int s =</mark>	0;	
Thread 1		Thread 2	
 compute f([A[i]) and put in reg0	9	 compute f([A[i]) and put in reg0	25
reg1 = s	0	reg1 = s	0
reg1 = reg1 + reg0	9	reg1 = reg1 + reg0	25
s = reg1		s = reg1	25
•••		•••	

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reg1 = s	0	reg1 = s	0
reg1 = reg1 + reg0	9	reg1 = reg1 + reg0	25
s = reg1	9	s = reg1	25
•••			

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A 3 5 f = square stat	<mark>ic int s =</mark>	0;	
Thread 1		Thread 2	
 compute f([A[i]) and put in reg0	9	 compute f([A[i]) and put in reg0	25
reg1 = s	0	reg1 = s	0
reg1 = reg1 + reg0	9	reg1 = reg1 + reg0	25
s = reg1	9	s = reg1	25
•••		•••	

- Assume A = [3,5], f is the square function, and s=0 initially
- For this program to work, s should be 34 at the end
  - but it may be 34, 9, or 25

#### Shared Memory code for computing a sum static int s = 0;

 static int s = 0;

 Thread 1

 local\_s1= 0

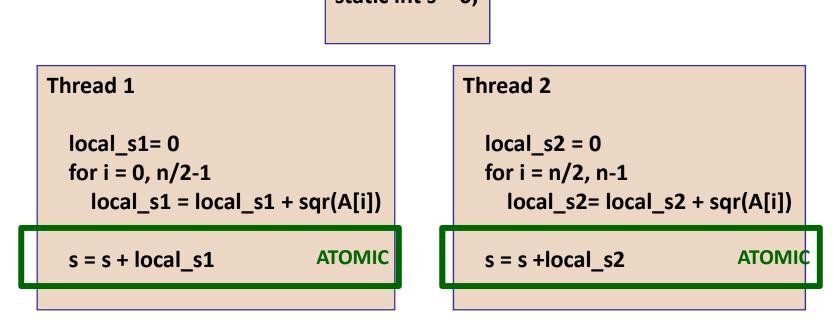
 for i = 0, n/2-1

 local\_s1 = local\_s1 + sqr(A[i])

 s = s + local\_s1

 s = s + local\_s1

#### Shared Memory code for computing a sum static int s = 0;



#### **Atomic Operations**

- To understand a concurrent program, we need to know what the **indivisible operations** are!
- Atomic Operation: an operation that always runs to completion or not at all
  - It is *indivisible*: it cannot be stopped in the middle and state cannot be modified by someone else in the middle
  - Fundamental building block if no atomic operations, then have no way for threads to work together
- On most machines, memory references and assignments (i.e. loads and stores) of words are atomic

#### Definitions

- Synchronization: using atomic operations to ensure cooperation between threads
  - For now, only loads and stores are atomic
  - hard to build anything useful with only reads and writes
- Mutual Exclusion: ensuring that only one thread does a particular thing at a time
  - One thread *excludes* the other while doing its task
- Critical Section: piece of code that only one thread can execute at once
  - Critical section and mutual exclusion are two ways of describing the same thing
  - Critical section defines sharing granularity

#### More Definitions

- Lock: prevents someone from doing something
  - Lock before entering *critical section* and before accessing *shared data*
  - Unlock when leaving, after accessing shared data
  - Wait if locked
    - Important idea: all synchronization involves waiting
- Example: fix the milk problem by putting a lock on refrigerator
  - Lock it and take key if you are going to go buy milk





#### Shared Memory code for computing a sum static int s = 0;

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 Thread 1

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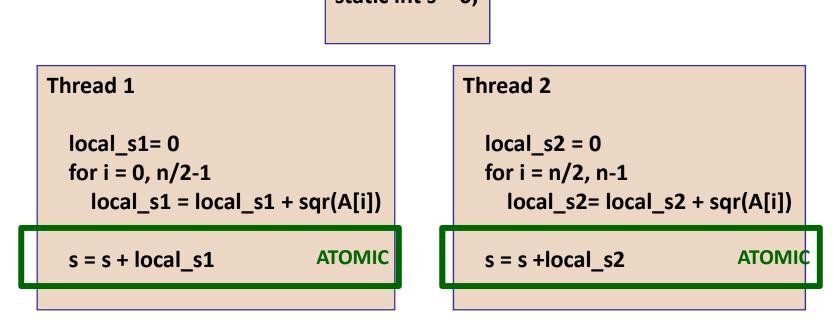
 for i = 0, n/2-1

 local\_s1 = local\_s1 + sqr(A[i])

 s = s + local\_s1

 s = s + local\_s1

#### Shared Memory code for computing a sum static int s = 0;



static int s = 0;
static lock lk;

Thread 1

```
local_s1= 0
for i = 0, n/2-1
    local_s1 = local_s1 + sqr(A[i])
    lock(lk);
s = s + local_s1
    unlock(lk);
```

#### Thread 2

local\_s2 = 0
for i = n/2, n-1
 local\_s2= local\_s2 + sqr(A[i])
lock(lk);
s = s +local\_s2
unlock(lk);

#### How to implement locks?

- Need HW support for **atomic** instructions
- RISCV uses two HW primitives
  - Load reserved
  - Store conditional
- (see discussion in *Chapter 2: Instructions, language of the computer,* slides 61-63)